



Value-Based Resource Management at SoC Scale

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Motivation

- Value-based schedulers have been demonstrated as an effective method in HPC domain
- Challenges in the SoC domain:
 - Dynamic workloads
 - Heterogeneous resource pool
 - Decision-making time
 - Effective utilization of heterogeneous SoC platforms
- First steps to adapt value-based scheduling heuristics to SoC scale

High Performance Computing

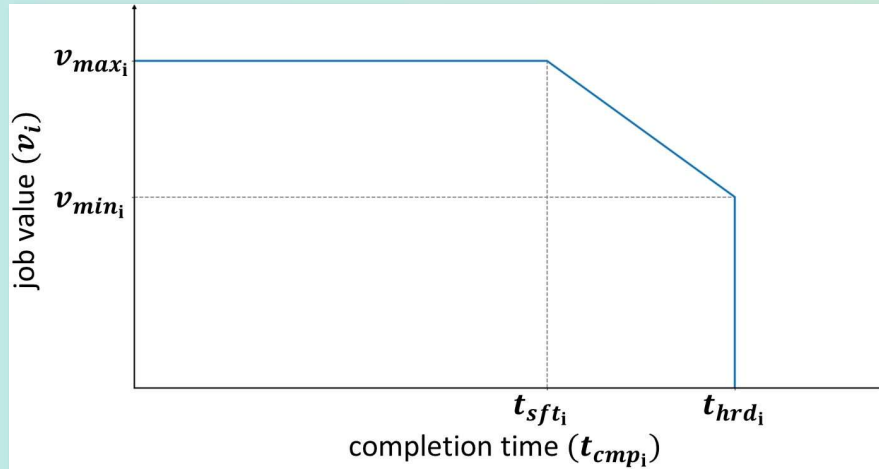


System on Chip



Value Definition

- Value of an application as a function of time:



$$v_i = \begin{cases} 0, & t_{cmp_i} > t_{hrd_i} \\ v_{max_i}, & t_{cmp_i} \leq t_{sft_i} \\ v_{max_i} - \frac{(t_{cmp_i} - t_{sft_i}) \times (v_{max_i} - v_{min_i})}{t_{hrd_i} - t_{sft_i}}, & t_{sft_i} < t_{cmp_i} < t_{hrd_i} \\ v_{min_i}, & t_{cmp_i} \geq t_{hrd_i} \end{cases}$$

Table 1: Definitions for Equation 1 and 2.

parameter	definition
i	application id in the order of arrival
$v_i(t)$	value earned by i^{th} application at time t
t_{hrd_i}	hard deadline for the application i
t_{sft_i}	soft deadline for the application i
t_{cmp_i}	completion time for the application i
v_{max_i}	maximum value application i can obtain
v_{min_i}	minimum value application i can obtain
t	time since the start of an experiment
$S(t)$	number of applications submitted till t
$SV(t)$	system value at time t

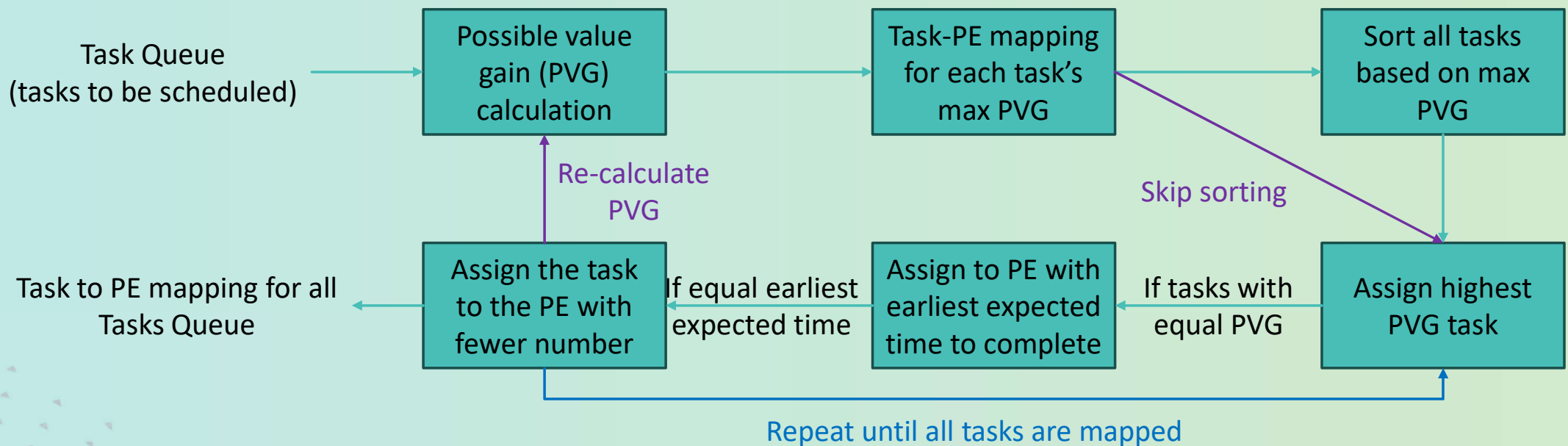
- Systems overall value gain:

$$\text{maximize } SV(t) = \sum_{i=1}^{S(t)} v_i(t_{cmp_i})$$

Value per Time Heuristic

- Two variations of value-based scheduling heuristic:
 - Value of Service (VoS)
 - Value of Service (VoS)

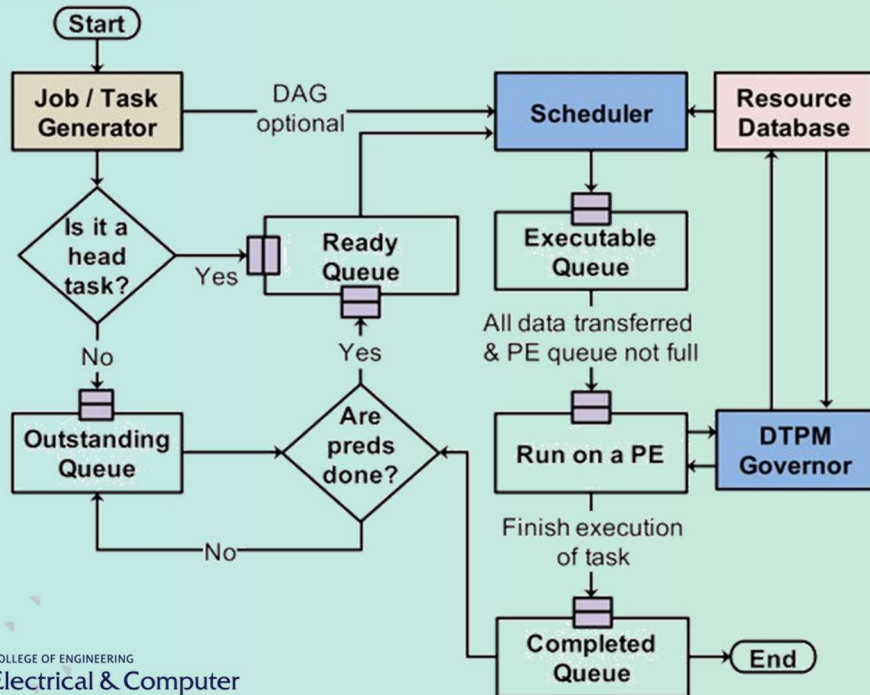
At any point, if PVG is calculated as 0 discard the task and mark the application as failed



Experiment Setup on DS3 and Applications

Simulation of Jetson AGX MPSoC on DS3¹:

- 7 CPUs that support all the tasks in the application DAGs
- 1 Accelerator that supports various tasks



Application Setup:

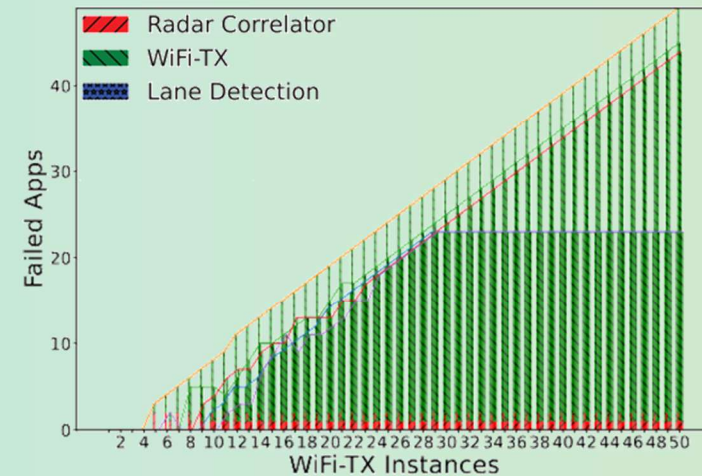
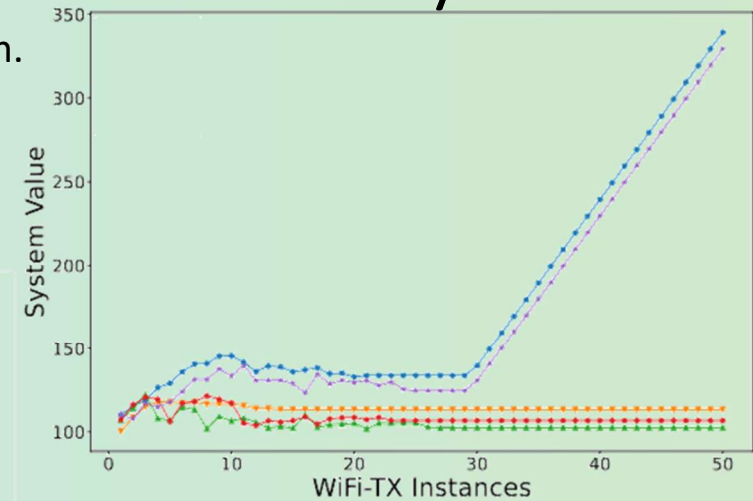
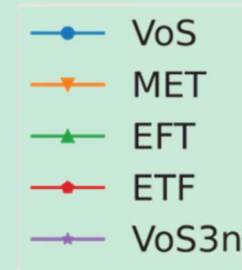
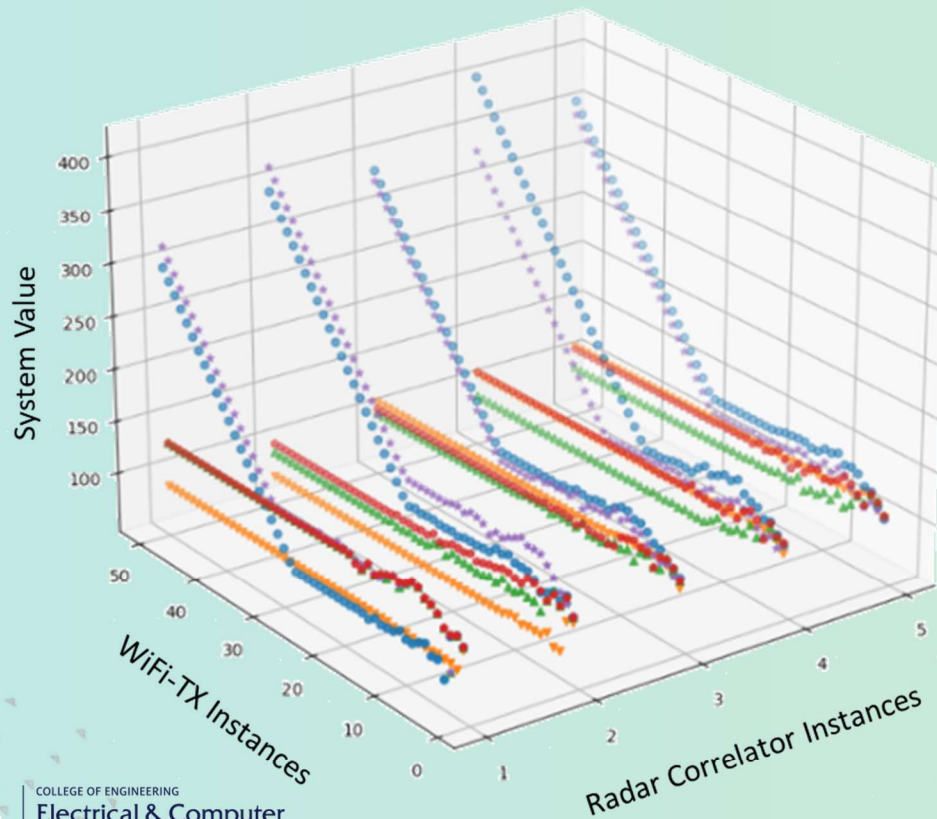
- Hard Threshold scaled as 2.1 ($t_{sft_i} = 2.1 \times EFT\ Exec$)
- Soft Threshold scaled as 1 ($t_{sft_i} = 1 \times EFT\ Exec$)
- v_{min_i} set to 0.5
- Injection rate set to $78.5\mu s$
 - One of each applications injected every $78.5\mu s$

	RC	WiFi-TX	LD
EFT Exec	5,865	61,910	78,522
v_{max}	1	10	100
Maximum Branches	2	5	1
Accelerateable Tasks	3	5	4
Acceleration Gain	10x	10x	15x-8x



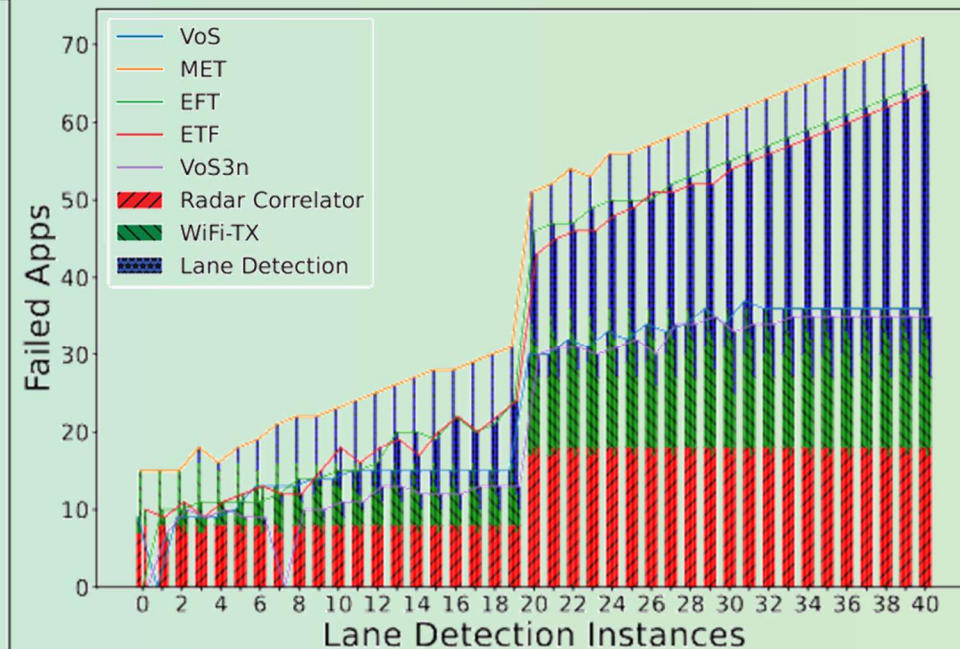
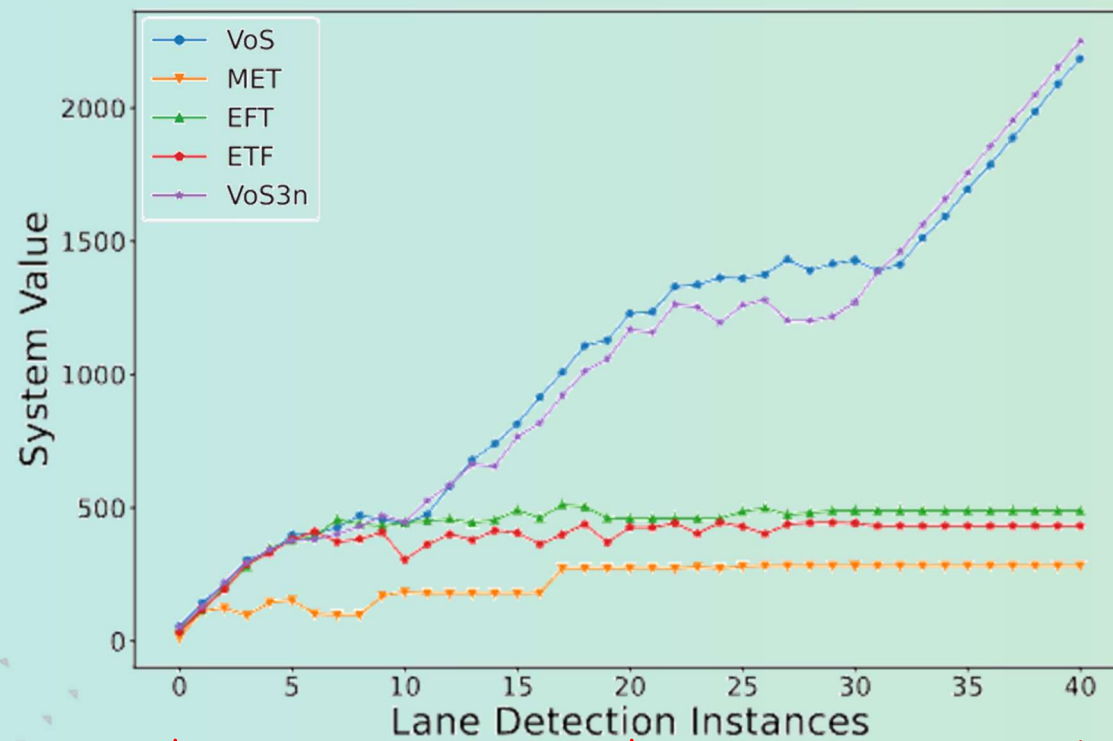
Results 1 – Undersubscribed to Oversubscribed System

- Value-based scheduler is more effective in application prioritization.



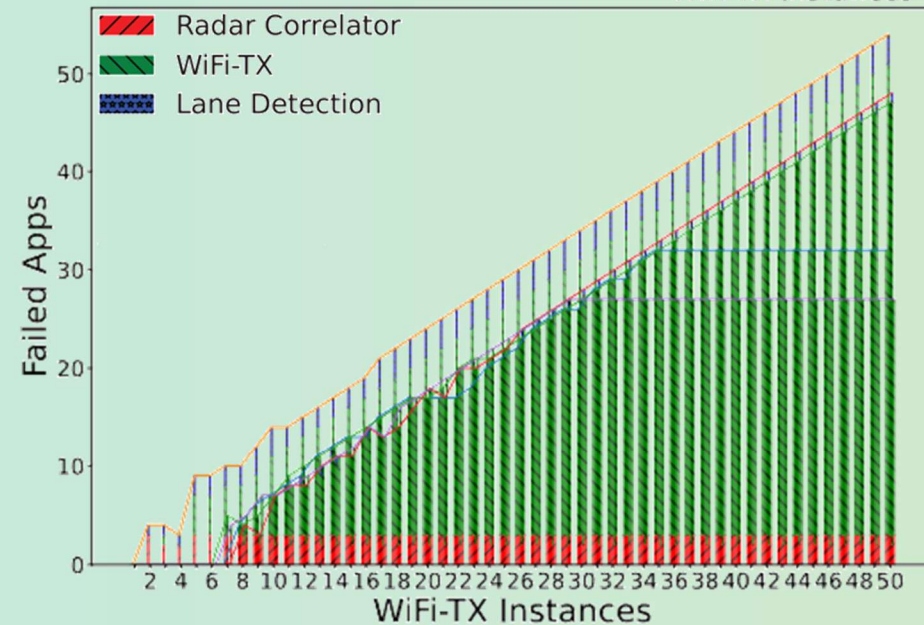
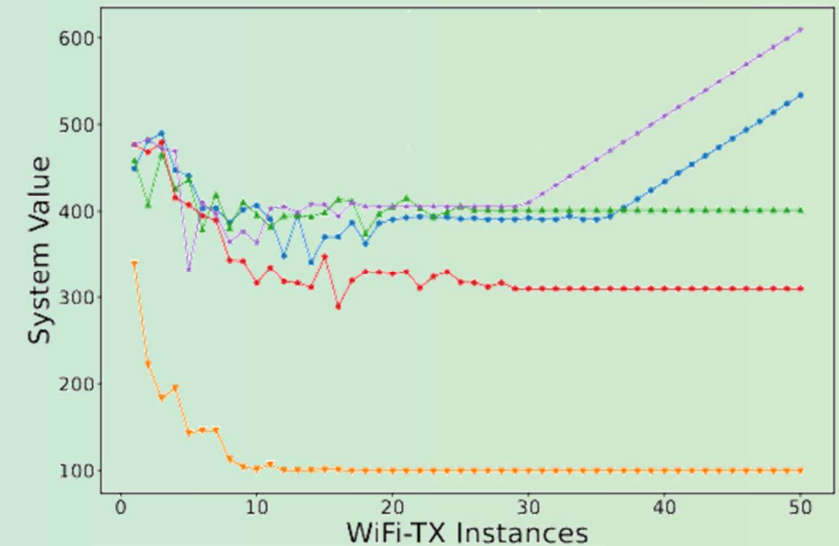
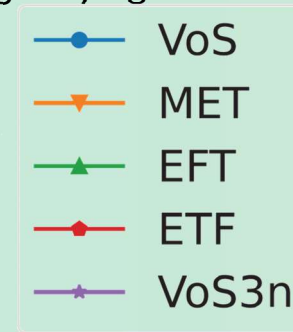
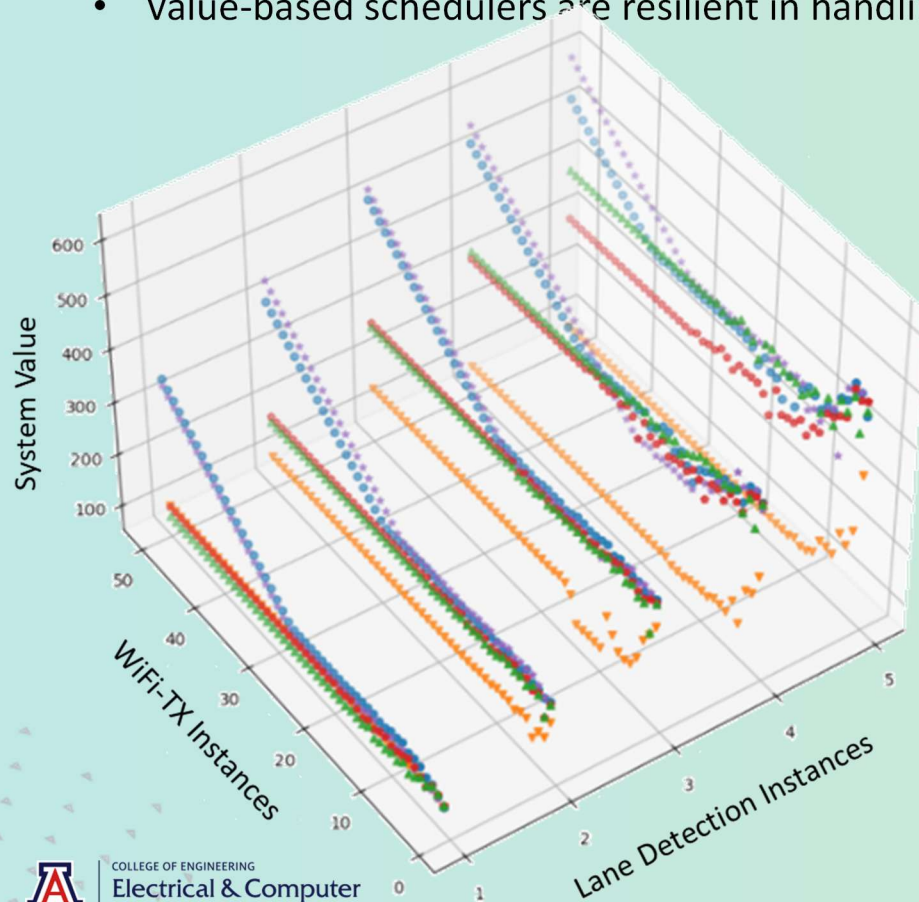
Results 2 – Reproducing Saturation Trends - Robustness and Adaptability

- Value-based schedulers are resilient in managing failed applications



Results 3 – Oversubscribed System

- Value-based schedulers are resilient in handling varying workloads



Conclusion and Future Work

- Use of value-based scheduling methods in heterogeneous SoC devices to maximize system productivity
- Value-based heuristics outperform traditional approaches
- Enabling the platforms to adapt dynamically to changing execution flows by incorporating task priority
- Dropping low-priority tasks in the decision-making process
- Stays in orders of $120\mu\text{s}$ scheduling overhead
- Moving toward dynamic priority changes



Thank You!

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Questions?



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